

Topology of $\mathcal{C} = 2^{\mathbb{N}}$

Lemma Every finitely determined set is **clopen**.

Pf: $U \in \mathcal{C}$ determined by $D \subset \mathbb{N}$ finite.

Open: $U \supseteq N_D(x)$ for each $x \in U$.

Closed: U^c is also fin. det. □

Thm: A set $U \in \mathcal{C}$ is clopen $\Leftrightarrow$ it is fin. det.

Pf: exercise!

Thm: The family of fin det sets forms a basis for the topology on $\mathcal{C}$.

Pf: Take $U \in \mathcal{C}$ open. Then $\forall x \in U$, there is finite $D(x) \subset \mathbb{N}$ s.t. $N_{D(x)}(x) \subseteq U$. Hence $U = \bigcup_{x \in U} N_{D(x)}(x)$. □

Lemma: The metric d on $\mathcal{C}$ is **complete** (every Cauchy seq. has a limit) ↪ exercise.

Summary:

① Top. on $\mathcal{C}$ is generated by a complete metric.

② $\mathcal{C}$ is **second-countable**: there is a countable basis for the topology

③ $\mathcal{C}$ is **zero-dimensional**: a basis of clopen sets

④ $\mathcal{C}$ is **compact** (exercise)

Prop: $\mathcal{B}(\mathcal{C}) =$ smallest σ -algebra on $\mathcal{C}$ containing all open sets.

Pf: exercise

→ Can define Borel sets in an arbitrary top. space.

Defn: A top. space X is Polish if:

- ① the topology is generated by a complete metric
- ② X is second-countable.

Examples: $\mathbb{C}$, $\mathbb{R}$, $\mathbb{R}^n$, $[0,1]$, $(0,1)$, $\mathbb{R} \setminus \mathbb{Q}$, ...
exercise!

Lemma: X Polish $\Rightarrow \exists$ a compatible complete metric bounded by 1.

Pf: $\tilde{d}(x,y) := \min\{d(x,y), 1\}$.

* Ctbl sums (i.e., disjoint unions) and products of Polish spaces are Polish.

Thm (Alexandrov): X a Polish space, $A \subseteq X$.
 A is Polish (in the subspace top) $\iff$
 A is G_δ .

ctbl $\cap$ of open sets

F_σ = ctbl $\cup$ of closed = complement of G_δ .

Why Polish spaces?

Thm X a Polish space $\Rightarrow$ there is a map code: $X \rightarrow \mathbb{C}$
s.t. $B \subseteq X$ is Borel in X iff $\{\text{code}(x) : x \in B\}$ is Borel in $\mathbb{C}$.

What is code?

Let $(U_n)_{n \in \mathbb{N}}$ be a basis for the top. on X .
if $x \in U_n$

Define $\text{code}(x) := (\text{code}_n(x))_{n \in \mathbb{N}}$ where $\text{code}_n(x) = \begin{cases} 0 & \text{if } x \notin U_n \\ 1 & \text{if } x \in U_n \end{cases}$.

Now, let $\Gamma := \text{graph}(\text{code}) \subseteq X \times \mathcal{C}$.

$$(x, y) \in \Gamma \iff y = \text{code}(x) \iff$$

$$\bigwedge_{n \in \mathbb{N}} \left(\underbrace{(y_n = 0 \ \& \ x \notin U_n)}_{\text{ctbl} \cap \text{Borel}(\text{clopen})} \vee \underbrace{(y_n = 1 \ \& \ x \in U_n)}_{\text{clopen} \cap \text{closed}} \right)$$

So Γ is a Borel set, hence code is a Borel function.

Since code is injective (exercise),

$$B \subseteq \mathcal{C} \text{ is Borel} \iff \text{code}(B) \text{ is Borel}$$

by Luzin-Suslin. □

More is true:

Thm (Borel isomorphism theorem):
 Let X be a Polish space, $A \subseteq X$ Borel, uncountable.
 Then $(A, \mathcal{B}(A)) \cong (\mathcal{C}, \mathcal{B}(\mathcal{C}))$

Defn: A standard Borel space is a structure isomorphic to $(X, \mathcal{B}(X))$ for a Polish space X .

↳ either $(X, \mathcal{P}(X))$ for ctbl X
or $\cong (\mathcal{C}, \mathcal{B}(\mathcal{C}))$.

Changing the topology

X a standard Borel space. A topology τ on X is compatible if $\mathcal{B}(X) = \text{Borel } \mathfrak{A}\text{-algebra of } (X, \tau)$.

Notation: $\langle \mathcal{U} \rangle = \text{top. generated by } \mathcal{U}$.

Thm: If $(\tau_n)_{n \in \mathbb{N}}$ are compatible Polish topologies on st. Borel X ,
then $\tau_\infty := \langle \bigcup_{n \in \mathbb{N}} \tau_n \rangle$ is also a compatible Polish top.

Pf: τ_∞ is clearly compatible, second-ctbl.

Fix complete metrics d_n bounded by 1 s.t. d_n generates τ_n .

$d_\infty(x, y) := \sum_{n \in \mathbb{N}} \frac{1}{2^{n+1}} d_n(x, y)$ is a complete metric that generates τ_∞ (exercise) ☺

Thm (MAKING BOREL SETS CLOPEN): X standard Borel.
 $\forall A \in \mathcal{B}(X)$, there is a compatible Polish τ on X
s.t. A is τ -clopen.

Pf: Let $\mathcal{A} := \{ A \in \mathcal{B}(X) : \dots \}$.

Fix a compatible Polish top. τ on X . Need to show:

- ① Every τ -open set is in $\mathcal{A}$
 - ② $\mathcal{A}$ is a σ -algebra
- $\Rightarrow \mathcal{A} \supseteq \mathcal{B}(X)$ ☺

①: $U \subseteq X$ τ -open $\Rightarrow U$ Polish, $X \setminus U$ Polish
open closed

$X \cong \underbrace{U \sqcup (X \setminus U)}_{\text{Polish}}, U \text{ and } X \setminus U \text{ both open. } \checkmark$

② $A \in \mathcal{A} \Rightarrow A^c \in \mathcal{A}$ clear

$(A_n \in \mathcal{A})_{n \in \mathbb{N}} \stackrel{?}{\Rightarrow} \bigcup A_n \in \mathcal{A}$

Let τ_n witness $A_n \in \mathcal{A}$

$\tau_\infty := \langle \bigcup_{n \in \mathbb{N}} \tau_n \rangle$ compatible, Polish

In τ_∞ , each A_n is clopen $\Rightarrow A$ is τ_∞ -open

$\Rightarrow A \in \mathcal{A}$ by ①.

→ Koll. / 2

