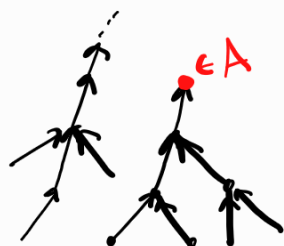


F is one-ended if there is no infinite F -descending sequence

$$x_0 \xleftarrow{F} x_1 \xleftarrow{F} x_2 \xleftarrow{F} x_3 \xleftarrow{\dots}$$

$G, A \subseteq V(G)$: A -one-ended subforest of G is a one-ended function

$F: A^c \rightarrow V(G)$ s.t. $F(x)$ is a nbr of x for all $x \in A^c$.



Proposition: G locally finite Borel graph, $A \subseteq V(G)$ a Borel set that meets every component of G , then there exists a Borel A -one-ended subforest.

Pf: $\text{dist}(x, A) =$ graph distance from x to A
 $=$ min # of edges in a path in G from x to A .

A path $x_0, x_1, \dots$ in G is A -repulsive if

$$\text{dist}(x_{n+1}, A) > \text{dist}(x_n, A) \quad \forall n.$$

$R := \{x \in A^c : \text{there is an infinite } A\text{-repulsive path starting at } x\}$

Last time: R is Borel. $\leftarrow$ this uses that G is locally finite

Note: $x \in R \Rightarrow x$ has a nbr $y \in R$ s.t. $\text{dist}(y, A) > \text{dist}(x, A)$

$x \in A^c \setminus R \Rightarrow x$ has a nbr y s.t. $\text{dist}(y, A) < \text{dist}(x, A)$

BIG idea #2: Borel choice for countable sets!

Thm: X, Y standard Borel, $A \subseteq X \times Y$ a Borel set.
 $\{x \in X : \exists! y \in Y \text{ s.t. } (x, y) \in A\}$ is Borel and $\neq \emptyset$

Suppose that for all $x \in X$, the fiber A_x is non-empty and closed.
 Then there is a Borel function $f: X \rightarrow Y$ s.t. $f(x) \in A_x \forall x \in X$.

f "chooses" a pt in A_x for each x

→ Proof: Apply Luzin-Novikov: Borel partial maps $f_n: X \dashrightarrow Y$
 s.t. $\forall x \in X, A_x = \{f_n(x) : x \in \text{dom}(f_n)\}$.

In particular, $\underline{f_n(x)} \in A_x$ for all $x \in \underline{\text{dom}(f_n)}$.

Define $N(x) := \min \{n \in \mathbb{N} : x \in \text{dom}(f_n)\}$ and

let $f(x) := f_{N(x)}(x)$. 😊

Using this, we obtain a Borel function $F: A^c \rightarrow V(\mathcal{U})$ s.t.:

$x \in R \Rightarrow F(x)$ is a nbr of x s.t. $F(x) \in R$ and
 $\text{dist}(F(x), A) > \text{dist}(x, A)$

$x \in A^c \setminus R \Rightarrow F(x)$ is a nbr of x s.t. $\text{dist}(F(x), A) < \text{dist}(x, A)$.

We claim that F is one-ended, as desired! 😊

Suppose not: $x_0, x_1, x_2, \dots$ is an F -descending sequence.

Case ①: All x_n are in R : $\Rightarrow \text{dist}(x_n, A) > \text{dist}(x_{n+1}, A)$
 for all $n \in \mathbb{N}$

$x_0 \xleftarrow{F} x_1 \xleftarrow{F} x_2 \xleftarrow{F} x_3 \xleftarrow{F} \dots$

$\text{dist}(x_0, A) > \text{dist}(x_1, A) > \text{dist}(x_2, A) > \dots$ **IMPOSSIBLE!**

Case ②: Some x_n is not in $R \Rightarrow x_m \notin R \forall m \geq n$, so

$x_n \xleftarrow{F} x_{n+1} \xleftarrow{F} x_{n+2} \xleftarrow{F} \dots$

$\notin R$

$\notin R$

$\notin R$

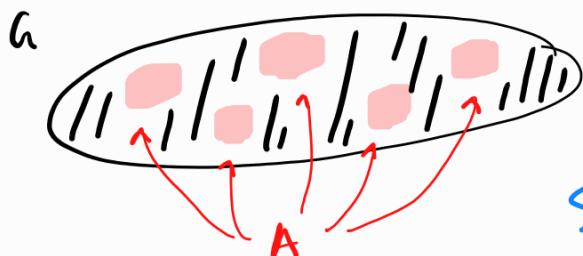
$\text{dist}(x_n, A) < \text{dist}(x_{n+1}, A)$

$< \text{dist}(x_{n+2}, A) < \dots$

$\Rightarrow x_n, x_{n+1}, x_{n+2}, \dots$ is an A-repulsive infinite path

$\Rightarrow x_n \in R$, a contradiction!

Corollary: G a Borel graph of max. deg. $d \in \mathbb{N}$,
 $A \subseteq V(G)$ a Borel set that meets every component of G ,
then $\chi_B(G-A) \leq d$.



Hope: Working with $G[A]$ is easier than working with G .

Specifically, we'll choose A so that every component of $G[A]$ is finite

BIG idea #3: Component-finite graphs are easy
(from the pt of view of Borel combinatorics)

Thm: Let G be a Borel component-finite graph.
Then $\chi_B(G) = \chi(G)$.

Proof: Let $V_n := \{x \in V(G) : \text{the component of } x \text{ has } n \text{ vertices}\}$
 $=$ the union of all components of size n

This is a Borel set by Luzin-Novikov

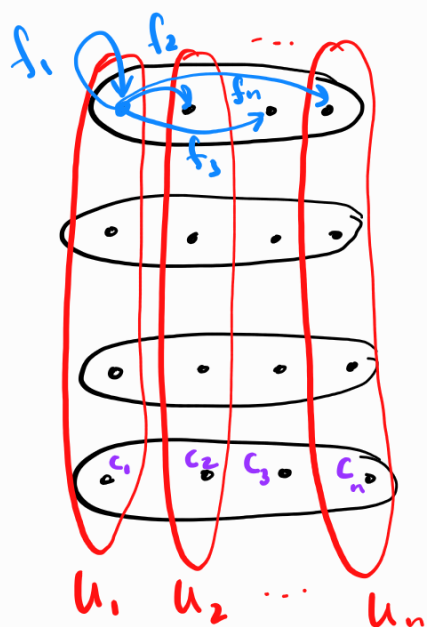
(exercise; hint: the relation " x and y are in the same component" is Borel)

$$V(G) = V_1 \cup V_2 \cup \dots$$

No edges between V_i and V_j for $i \neq j$, so we may color each V_i separately.

So fix $n \in \mathbb{N}$ and pass to $G[V_n]$

→ from now on, assume that every component of h has exactly n vertices.



Let $<$ be a Borel linear order on $V(h)$. For $1 \leq i \leq n$, let

$U_i := \{x : x \text{ is the } i^{\text{th}} \text{ vtx of its component according to } <\}$

↑
Borel by LN

Also, let $f_i(x)$ be the i^{th} vtx in the component of x

(i.e., $f_i(x)$ is the unique vtx from U_i in the same component as x)

Let $k := \chi(h)$. Want: find a Borel k -coloring.

$A := \{ (x, c) \in U_1 \times \{0, \dots, k-1\}^n :$

$\forall 1 \leq i < j \leq n : \{f_i(x), f_j(x)\} \in E(h) \Rightarrow c_i \neq c_j \}$

For each $x \in U_1$, the fiber A_x is finite and $\neq \emptyset$ $\nearrow \chi(h) = k$

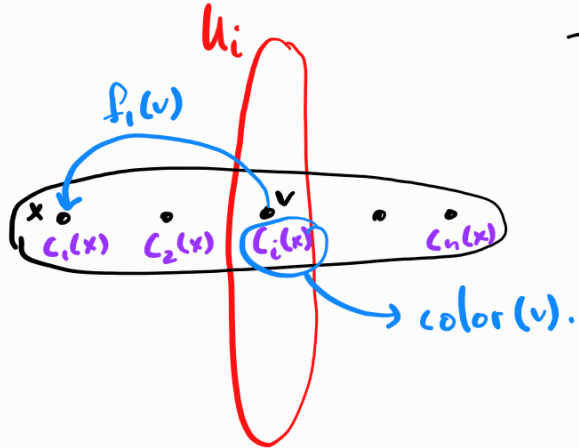
$\Rightarrow$ there is a Borel function $c: U_1 \rightarrow \{0, \dots, k-1\}^n$

s.t. $c(x) \in A_x$ for all $x \in U_1$.

$c = (c_1, \dots, c_n), c_i: U_1 \rightarrow \{0, \dots, k-1\}$

Define, for $v \in V(h)$,

$\text{color}(v) := c_i(f_i(v))$ where $v \in U_i$.



This is a proper Borel k -coloring, as desired! 🧠



How to extend the coloring to A ?

Different sets of available colors to different vertices...

BIG idea #4: Use fancy results from finite combinatorics/graph theory! 😊